

BTCWare Ransomware

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Last years' news headlines were dominated by ransomware attacks like Wannacry and Petya, the constant presence of Cerber, the disappearance and return of Locky, and the growing popularity of Ransomware-as-a-Service (RaaS) and smaller campaigns like Jaff and BadRabbit.

Introduction

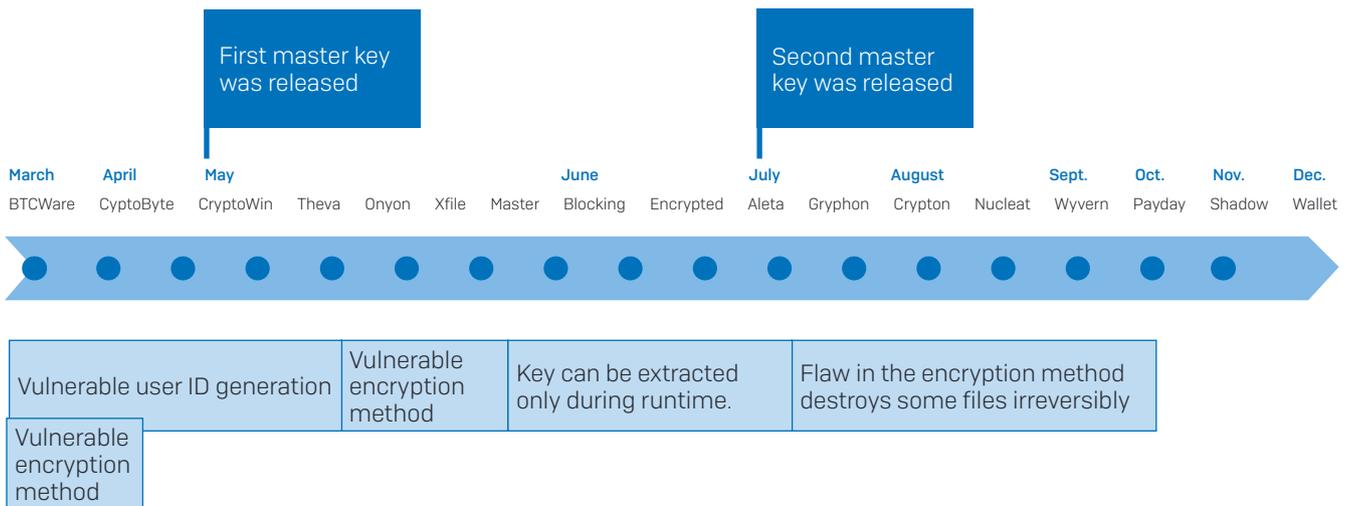
But there are many ransomware families that didn't get much attention, even though they were also active. One such example is BTCWare, which was among the 10 most commonly intercepted ransomware families of the year.

We've collected and analyzed 17 different variants of this family, from its first appearance in March 2017 until today. While looking into these variants, we can follow the development process of the BTCWare ransomware family and see how the developers have made mistakes, learned from them, and tried to improve their code over time.

The authors of this family experimented with different key generation and encryption methods. Looking through these solutions, we will introduce a few implementation failures and the route which led from a completely vulnerable method to a cryptographically secure file encryption. Although the analysis of the file encryption method will be a focus of this paper, we will cover some other features of the BTCWare family and highlight the similarities and differences of different variants.

We refer to the variants by the new extension they add to the encrypted file names, but it is important to note that there might be differences within variants as well. It is possible that, for example, the change in the encryption method didn't happen at the same time as the change in the extension.

On the timeline below, we have summarized the appearance of the different variants by month:



Infection vector

The BTCWare ransomware family targets Windows machines and is primarily distributed by brute-forcing weak passwords of the Remote Desktop Protocol (RDP) and [manually installing the malware](#).

A few variants of BTCWare were distributed by spam emails as well. The so-called “Blank Slate” malspam campaign is responsible for spreading different ransomware families targeting Windows [e.g. Cerber, Locky, Globelmposter] ¹.

In a few cases, BTCWare variants were distributed via this campaign. ^{2,3} We observed a big spike in August that spread the Gryphon variant. Previously, in July, Aleta was distributed using this method.

These spam emails arrive with no subject and no content. The attachment is either a Zip file containing another Zip archive with a JavaScript file inside or a Microsoft Word document. Once JavaScript or a malicious macro is executed, the BTCWare payload is downloaded.

Execution

Apart from continuous updates to the file encryption method, BTCWare performs additional changes to the system such as:

- Adding an auto run entry in the registry under Run key value,
- Deleting volume shadow copies using a vssadmin tool,
- Disabling recovery settings using a bcdedit tool.

Below is the summary of various system changes done by different BTCWare variants:

Variant	Ransom note name	Added to HKCU\Software\Microsoft\Windows\CurrentVersion\Run registry key	Mutex name
BTCWare	#_HOW_TO_FIX.inf	Value Name: btq Value Data: %Appdata%\mfskSkfkl.exe (this is a copy of the original executable) Deleted after encryption is done.	BTCWXXX
Cryptobyte	#_HOW_TO_FIX.inf	Value Name: btq Value Data: %Appdata%\mfskSkfkl.exe (this is a copy of the original executable) Deleted after encryption is done. Value Name: insbtq Value Data: %Appdata%\#_HOW_TO_FIX.inf	BTCWXXX
Cryptowin	#_HOW_TO_FIX.inf	Value Name: btq Value Data: %Appdata%\mfskSkfkl.exe (this is a copy of the original executable) Deleted after encryption is done. Value Name: insbtq Value Data: %Appdata%\#_HOW_TO_FIX.inf	BTCWXXX
Theva	#_README_#.inf	Value Name: xcz Value Data: %Appdata%\zmsksdffff.exe (this is a copy of the original executable) Deleted after encryption is done. Value Name: info Value Data: %Appdata%\#_README_#.inf	CRYNXXX
Onyon	!#_DECRYPT_#!.inf	Value Name: DECRYPTINFO Value Data: %Appdata%\!#_DECRYPT_#!.inf	ONYONLOCK
Xfile	!#_DECRYPT_#!.inf	Value Name: DECRYPTINFO Value Data: %Appdata%\!#_DECRYPT_#!.inf	ONYONLOCK
Master	!#_RESTORE_FILES_#!.inf	Value Name: DECRYPTINFO Value Data: %Appdata%\!#_RESTORE_FILES_#!.inf	ONYONLOCK or MASTERLOCK

Encrypted	!#_RESTORE_FILES_#!.inf	Value Name: DECRYPTINFO Value Data: %Appdata%!#_RESTORE_FILES_#!.inf	MASTERLOCK
Blocking	!#_RESTORE_FILES_#!.inf	Value Name: DECRYPTINFO Value Data: %Appdata%!#_RESTORE_FILES_#!.inf	MASTERLOCK
Aleta	!#_READ_ME_#!.inf !#_READ_ME_#!.hta %Appdata%\Info.hta	Value Name: DECRYPTINFO Value Data: %Appdata%!#_READ_ME_#!.inf Value Name: DECRYPTINFO Value Data: %Appdata%\ Info.hta	MASTERLOCK
Gryphon	!## DECRYPT FILES##!.txt %Appdata%\HELP.txt	Value Name: DECRYPTINFO Value Data: %Appdata%\HELP.txt	no mutex or GIVEMEBTC
Crypton	!## DECRYPT FILES##!.txt %Appdata%\HELP.txt	Value Name: DECRYPTINFO Value Data: %Appdata%\HELP.txt	no mutex or GIVEMEBTC
Nuclear	HELP.hta	Value Name: DECRYPTINFO Value Data: %Appdata%\HELP.hta Value Name: HELPINFO Value Data: %Appdata%\HELP.hta	NUCLEAR
Wyvern	!FILES ENCRYPTED.txt %Appdata%\HELP.hta	Value Name: DECRYPTINFO Value Data: %Appdata%\HELP.hta	NUCLEAR
Payday	!! RETURN FILES !!.txt %Appdata%\payday.hta	Value Name: payday Value Data: %Appdata%\payday.hta Value Name: baby Value Data: %Appdata%\payday.hta	PAYDAYDAYPAY
Shadow	!! RETURN FILES !!.txt %Appdata%\payday.hta	Value Name: 1payday Value Data: %Appdata%\payday.hta Value Name: 2baby Value Data: %Appdata%\payday.hta	PAYDAYDAYPAY
Wallet	! How Decrypt Files.txt %Appdata%\payday.hta	Value Name: 1payday Value Data: %Appdata%\payday.hta Value Name: 2baby Value Data: %Appdata%\payday.hta	PAYDAYDAYPAY

Table 1

When BTCWare is executed, it starts by checking a specific mutex. If the mutex exists, the process ends. Otherwise it creates the mutex. The different mutex names can be found in Table 1.

In the case of the first few variants, it creates a copy of itself in the %Appdata% folder and adds this to the HKCU\Software\Microsoft\Windows\CurrentVersion\Run registry key to ensure it runs after reboot. After the file encryption is completed, the created registry values are deleted.

All the variants save the ransom note in the %Appdata% folder and add this to the HKCU\Software\Microsoft\Windows\CurrentVersion\Run registry key. This way, every time the infected computer is rebooted, the ransom note will be opened. From the last two variants (Shadow and Wallet), this information is added to the HKLM\Software\Microsoft\Windows\CurrentVersion\Run registry key as described in the above Table 1.

To reduce the chance of recovering the encrypted files, BTCWare attempts to delete volume shadow copies using the vssadmin.exe tool with the following command:

vssadmin.exe Delete Shadows /All /Quiet

Using bcdedit.exe, BTCWare disables recovery and boot failures using the two commands below:

bcdedit.exe /set {default} bootstatuspolicy ignoreallfailures

bcdedit.exe /set {default} recoveryenabled No

```

00000000403AC3
00000000403AC3
00000000403AC3 8B 35 DC 81 42 00
00000000403AC9 6A 00
00000000403ACB 6A 00
00000000403ACD 68 48 52 43 00
00000000403AD2 68 A0 52 43 00
00000000403AD7 68 B0 52 43 00
00000000403ADC 6A 00
00000000403ADE FF D6
00000000403AE0 6A 00
00000000403AE2 6A 00
00000000403AE4 68 C0 52 43 00
00000000403AE9 68 A0 52 43 00
00000000403AEE 68 B0 52 43 00
00000000403AF3 6A 00
00000000403AF5 FF D6
00000000403AF7 6A 00
00000000403AF9 6A 00
00000000403AFB 68 28 53 43 00
00000000403B00 68 A0 52 43 00
00000000403B05 68 B0 52 43 00
00000000403B0A 6A 00
00000000403B0C FF D6
00000000403B0E 6A 00
00000000403B10 6A 00
00000000403B12 68 B0 53 43 00
00000000403B17 68 A0 52 43 00
00000000403B1C 68 B0 52 43 00
00000000403B21 6A 00
00000000403B23 FF D6

loc_403AC3:
mov     esi, ds:ShellExecuteW
push   0           ; nShowCmd
push   0           ; lpDirectory
push   offset Parameters ; "/c vssadmin.exe Delete Shadows /All /Qu"...
push   offset File      ; "cmd.exe"
push   offset Operation ; "open"
push   0           ; hwnd
call   esi ; ShellExecuteW
push   0           ; nShowCmd
push   0           ; lpDirectory
push   offset aCBcdeditExeSet ; "/c bcdedit.exe /set {default} recoverye"...
push   offset File      ; "cmd.exe"
push   offset Operation ; "open"
push   0           ; hwnd
call   esi ; ShellExecuteW
push   0           ; nShowCmd
push   0           ; lpDirectory
push   offset aCBcdeditExeSet_0 ; "/c bcdedit.exe /set {default} bootstatu"...
push   offset File      ; "cmd.exe"
push   offset Operation ; "open"
push   0           ; hwnd
call   esi ; ShellExecuteW
push   0           ; nShowCmd
push   0           ; lpDirectory
push   offset aCVssadminExeDe_0 ; "/c vssadmin.exe delete shadows /all /qu"...
push   offset File      ; "cmd.exe"
push   offset Operation ; "open"
push   0           ; hwnd
call   esi ; ShellExecuteW
    
```

In the first few variants, the strings related to the above mentioned registry keys, commands and ransom note were not obfuscated or encrypted at all in the binary.

With the Aleta variant, the ransom note was stored base64 encoded. Later all the above-mentioned strings were base64 encoded as well: from Nuclear the email address, from Payday all the strings related to the registry entries and the commands above.

Encryption

BTCWare kept evolving by making constant updates to its encryption method.

In Table 2, we have summarized a few details of the file encryption method of the different BTCWare variants.

Variant	New extension	User ID	File encryption algorithm	Encrypted size	Leaked master key	Where is the encrypted user ID stored?
BTCWare	.[< email address >].btcware	<0x19 random alphanumeric characters>	RC4	Whole file in 1000 byte long chunks	yes	key.dat
Cryptobyte	[< email address >].cryptobyte	OBAMA-<0x19 random alphanumeric characters>-<YYYY>-<MM>-<DD> XZD-<0x19 random alphanumeric characters>-<YYYY>-<MM>-<DD>	RC4	Whole file in 1000 byte long chunks	yes	Ransom note
Cryptowin	.[< email address >].cryptowin	ADM-<0x19 random alphanumeric characters>-<YYYY>-<MM>-<DD>	AES-192	Whole file in 992 byte long chunks	yes	Ransom note
Theva	.[< email address >].theva	SX-<random alphanumeric string>-<YYYY>-<MM>-<DD>	AES-192	Whole file in 992 byte long chunks	yes	Ransom note
Onyon	.[< email address >].onyon or .onyon	SX/DN-<random alphanumeric string>-<YYYY>-<MM>-<DD> or OBAMA-<random alphanumeric string>-<YYYY>-<MM>-<DD>	RC4	Encrypted length max 0xA00000 bytes	yes	Ransom note
Xfile	.[< email address >].xfile or .xfile	SX-<random alphanumeric string>-<YYYY>-<MM>-<DD>	RC4	Encrypted length max 0xA00000 bytes	yes	Ransom note

Master	.[< email address >].master	<number>-<random alphanumeric string>-<YYYY>-<MM>-<DD>	AES-256	Encrypted length max 0xA00000 bytes	yes	Ransom note
Encrypted	.[< email address >].encrypted	<number>-<random alphanumeric string>-<YYYY>-<MM>-<DD>	AES-256	Encrypted length max 0xA00000 bytes	no	Ransom note
Aleta	.[< email address >].aleta	<number>-<random alphanumeric string>-<YYYY>-<MM>-<DD>	AES-256	Encrypted length max 0xA00000 bytes	no	Ransom note or key.aleta
Blocking	.[< email address >].blocking	<number>-<random alphanumeric string>-<YYYY>-<MM>-<DD>	AES-256	Encrypted length max 0xA00000 bytes	no	Ransom note
Gryphon	.[< email address >].gryphon	User ID is generated using function addresses and the result of GetTickCount	AES-256	Encrypt first 0xFA00 bytes	no	Ransom note
Crypton	.[< email address >].crypton	User ID is generated using function addresses and the result of GetTickCount	AES-256	Encrypt first 0xFA00 bytes	no	Ransom note
Nuclear	.[< email address >].nuclear	User ID is generated using function addresses and the result of GetTickCount	AES-256	Encrypted length max 0xA00000 bytes	no	At the end of the encrypted file
Wyvern	.[< email address >-id-< id >].wyvern	User ID is generated using function addresses and the result of GetTickCount	AES-256	Encrypted length max 0xA00000 bytes	no	At the end of the encrypted file
Payday	.[< email address >-id-< id >].payday	CryptGenRandom is used for key generation	AES-256	Encrypted length max 0xA00000 bytes	no	At the end of the encrypted file
Shadow	.[< email address >-id-< id >].shadow	CryptGenRandom is used for key generation	AES-256	Encrypted length max 0xA00000 bytes	no	At the end of the encrypted file
Wallet	.[< email address >-id-< id >].wallet or .[< email address > or < email address >-id-< id >].wallet	CryptGenRandom is used for key generation	AES-256	Encrypted length max 0xA00000 bytes	no	At the end of the encrypted file

Table 2

In each case, a new extension is appended to the encrypted file name. This extension usually contains an email address, which can be used to contact the attackers. Although the details of the key generation and the used encryption algorithm vary over different variants, the general process of file encryption is very similar in most variants.

BTCWare uses the combination of asymmetric and symmetric encryption. Each sample contains a hard-coded RSA-1024 public key, which means that this ransomware can encrypt files offline without communicating with a command-and-control server.

The file encryption process can be summarized in four steps:

1. User ID generation
2. User ID encryption
3. Symmetric key derivation
4. File encryption

In the next four sections, we will examine these steps and highlight the similarities and differences between different variants.

User ID generation

The first step is to generate a user ID, which later will be used for symmetric key generation. The list of the user IDs can be found in the third column of Table 2. We can group the variants into six different categories based on the used method to generate the user ID. Let's look at these one by one.

BTCWare

The first variant generates 0x19 random alphanumeric characters using the `__time64`, `_srand` and `_rand` functions. The first call to `_rand` determines if a number, lowercase letter or uppercase letter will be generated and the next call to `_rand` generates the character.

Example: 1HYvsz449Uhn4gs68N1vy0U5a

The return value of `__time64` in the EAX register is used as seed for the pseudorandom number generation. This is a 32 bit value, which can be easily brute-forced. This and the fact that the same key was used to encrypt all the files make it possible to decrypt the encrypted files without paying the ransom.

```
.text:004036F0 generate_rand proc near ; CODE XREF: WinMain(x,x,x,x)+FF4p
.text:004036F0 push ebx
.text:004036F1 push esi
.text:004036F2 push edi
.text:004036F3 push 0 ; Time
.text:004036F5 call __time64
.text:004036FA push eax ; unsigned int
.text:004036FB call _srand
.text:00403700 push 1Ah ; unsigned int
.text:00403702 call ??_U@YAPAXI@Z ; operator new[](uint)
.text:00403707 add esp, 0Ch
.text:0040370A mov edi, eax
.text:0040370C xor esi, esi
.text:0040370E lea ebx, [esi+1Ah]
.text:00403711
.text:00403711 loc_403711: ; CODE XREF: generate_rand+6E4j
.text:00403711 call _rand
.text:00403716 cdq
.text:00403717 mov ecx, 3
.text:0040371C idiv ecx
.text:0040371E sub edx, 0
.text:00403721 jz short loc_403747
.text:00403723 sub edx, 1
.text:00403726 jz short loc_40373A
.text:00403728 sub edx, 1
.text:0040372B jnz short loc_40375A
.text:0040372D call _rand
.text:00403732 cdq
.text:00403733 idiv ebx
.text:00403735 add dl, 61h ; 'a' ; lowercase letters
.text:00403738 jmp short loc_403757
.text:0040373A ; -----
.text:0040373A loc_40373A: ; CODE XREF: generate_rand+361j
.text:0040373A call _rand
.text:0040373F cdq
.text:00403740 idiv ebx
.text:00403742 add dl, 41h ; 'A' ; uppercase letters
.text:00403745 jmp short loc_403757
.text:00403747 ; -----
.text:00403747 loc_403747: ; CODE XREF: generate_rand+311j
.text:00403747 call _rand
.text:0040374C cdq
.text:0040374D mov ecx, 0Ah
.text:00403752 idiv ecx
.text:00403754 add dl, 30h ; '0' ; number 0-9
```

Figure 1: BTCWare random generation

Cryptobyte, Cryptowin, Theva

0x19 random values are generated the same way as previously, but other characters are added to them.

Examples:

XZD-1ZCgy2V1tLUA22weX2PwOK9lc-2018-01-11

ADM-5i051f4066j6JaK04JA8qfLt6-2018-01-11

OBAMA-GiaGHJ03PTyN07J9V011F5pH5-2018-01-11

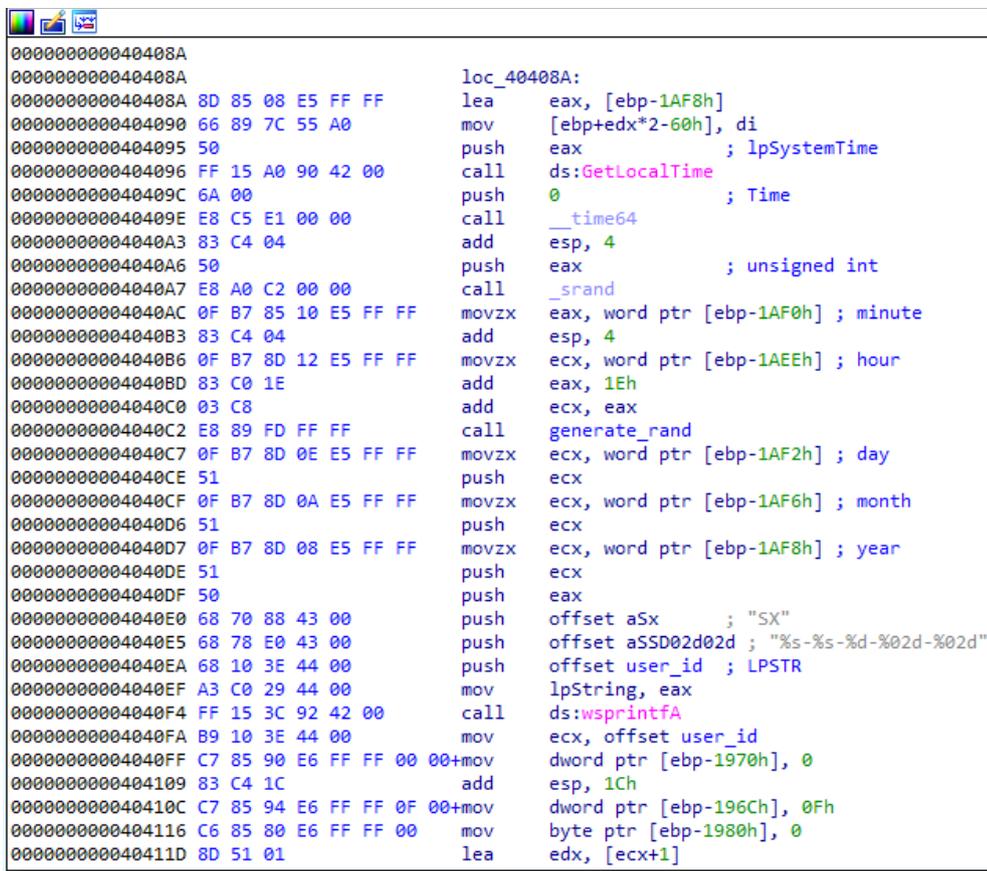
In case of Theva variant, the number of random alphanumeric characters is calculated using the hour and minute value from the result of GetLocalTime function.

Example:

SX-IMJ93SpGA01sr29UK04M6f5TziTbjgQtH2R03358oh4k8Kf3YYt7BTL3Gcy7WD567F0m-2018-01-11.

Only the random alphanumeric characters will be used to derive the file encryption key.

The Cryptobyte and Cryptowin user IDs can be brute-forced the same way as in case of the BTCWare variant. Meanwhile the changes in Theva only multiply the number of possible keys by 1440, which means that brute-forcing is still possible.



```

0000000040408A      loc_40408A:
0000000040408A      lea     eax, [ebp-1AF8h]
0000000040408A      8D 85 08 E5 FF FF   mov     [ebp+edx*2-60h], di
00000000404090      66 89 7C 55 A0     mov     eax, di
00000000404095      50             push   eax             ; lpSystemTime
00000000404096      FF 15 A0 90 42 00   call   ds:GetLocalTime
0000000040409C      6A 00             push   0             ; Time
0000000040409E      E8 C5 E1 00 00     call   __time64
000000004040A3      83 C4 04             add     esp, 4
000000004040A6      50             push   eax             ; unsigned int
000000004040A7      E8 A0 C2 00 00     call   _srand
000000004040AC      0F B7 85 10 E5 FF FF   movzx  eax, word ptr [ebp-1AF0h] ; minute
000000004040B3      83 C4 04             add     esp, 4
000000004040B6      0F B7 8D 12 E5 FF FF   movzx  ecx, word ptr [ebp-1AEEh] ; hour
000000004040BD      83 C0 1E             add     eax, 1Eh
000000004040C0      03 C8             add     ecx, eax
000000004040C2      E8 89 FD FF FF     call   generate_rand
000000004040C7      0F B7 8D 0E E5 FF FF   movzx  ecx, word ptr [ebp-1AF2h] ; day
000000004040CE      51             push   ecx
000000004040CF      0F B7 8D 0A E5 FF FF   movzx  ecx, word ptr [ebp-1AF6h] ; month
000000004040D6      51             push   ecx
000000004040D7      0F B7 8D 08 E5 FF FF   movzx  ecx, word ptr [ebp-1AF8h] ; year
000000004040DE      51             push   ecx
000000004040DF      50             push   eax
000000004040E0      68 70 88 43 00     push   offset aSx      ; "SX"
000000004040E5      68 78 E0 43 00     push   offset aSSD02d02d ; "%s-%s-%d-%02d-%02d"
000000004040EA      68 10 3E 44 00     push   offset user_id  ; LPSTR
000000004040EF      A3 C0 29 44 00     mov     lpString, eax
000000004040F4      FF 15 3C 92 42 00   call   ds:wprintfA
000000004040FA      B9 10 3E 44 00     mov     ecx, offset user_id
000000004040FF      C7 85 90 E6 FF FF 00 00+mov  dword ptr [ebp-1970h], 0
00000000404109      83 C4 1C             add     esp, 1Ch
0000000040410C      C7 85 94 E6 FF FF 0F 00+mov  dword ptr [ebp-196Ch], 0Fh
00000000404116      C6 85 80 E6 FF FF 00   mov     byte ptr [ebp-1980h], 0
0000000040411D      8D 51 01             lea     edx, [ecx+1]

```

Figure 2: Theva user ID generation (zmsksddfff.exe)

Onyon, Xfile

To create the user ID, these variants first generate a random string [string1] using GetLocalTime, the address of a few functions and a similar random generation method as previously. Then a string like "SX" or "DN" is added to the front of the user ID and the date to the end.

Example: DN-1967722968a38qpA50x6avM3TRT7lHe3xm82bKW913712A19677236001966031880YSJy9752FCEE9-2018-01-11

The generated random string [string1] is used to derive the file encryption key. Adding the address of certain functions is an unusual way to increase randomness, but this, with using the random generation twice, ensures that the previous brute-force method cannot be used anymore.

```

000000004035CA
000000004035CA loc_4035CA:
000000004035CA FF 35 EC 71 42 00 push ds:GetClipboardOwner
000000004035D0 51 push ecx
000000004035D1 B9 05 00 00 00 mov ecx, 5
000000004035D6 E8 A5 F7 FF FF call generate_rand
000000004035DB 8B 8D C8 F1 FF FF mov ecx, [ebp-0E38h]
000000004035E1 83 C4 04 add esp, 4
000000004035E4 50 push eax
000000004035E5 FF 35 E0 71 42 00 push ds:GetActiveWindow
000000004035EB 0F B7 85 D8 F1 FF FF movzx eax, word ptr [ebp-0E28h] ; second
000000004035F2 FF 35 48 70 42 00 push ds:GetCurrentThreadId
000000004035F8 50 push eax
000000004035F9 0F B7 85 DA F1 FF FF movzx eax, word ptr [ebp-0E26h] ; millisecond
00000000403600 50 push eax
00000000403601 0F B7 85 D4 F1 FF FF movzx eax, word ptr [ebp-0E2Ch] ; hour
00000000403608 83 C0 05 add eax, 5
0000000040360B 03 C8 add ecx, eax
0000000040360D E8 6E F7 FF FF call generate_rand
00000000403612 50 push eax
00000000403613 FF 35 60 70 42 00 push ds:GetCurrentProcessId
00000000403619 68 64 38 43 00 push offset aSDXDDSX ; "%d%s%d%0Xd%d%s%X"
0000000040361E FF 35 70 DC E3 00 push string1 ; LPSTR
00000000403624 FF 15 E8 71 42 00 call ds:wsprintfA
0000000040362A 0F B7 85 D2 F1 FF FF movzx eax, word ptr [ebp-0E2Eh] ; day
00000000403631 50 push eax
00000000403632 0F B7 85 CE F1 FF FF movzx eax, word ptr [ebp-0E32h] ; month
00000000403639 50 push eax
0000000040363A 0F B7 85 CC F1 FF FF movzx eax, word ptr [ebp-0E34h] ; year
00000000403641 50 push eax
00000000403642 FF 35 70 DC E3 00 push string1
00000000403648 68 E0 37 43 00 push offset aSx ; "SX"
0000000040364D 68 78 38 43 00 push offset aSSD02d02d ; "%s-%s-%d-%02d-%02d"
00000000403652 68 48 B5 E3 00 push offset user_id ; LPSTR
00000000403657 FF 15 E8 71 42 00 call ds:wsprintfA
0000000040365D B9 48 B5 E3 00 mov ecx, offset user_id
00000000403662 C7 85 C0 F1 FF FF 00 00 mov dword ptr [ebp-0E40h], 0
0000000040366C 83 C4 44 add esp, 44h
0000000040366F C7 85 C4 F1 FF FF 0F 00 mov dword ptr [ebp-0E3Ch], 0Fh
00000000403679 C6 85 B0 F1 FF FF 00 mov byte ptr [ebp-0E50h], 0
00000000403680 8D 51 01 lea edx, [ecx+1]

```

Figure 3: Onyon user ID generation

Even though the user-ID generation flaw is fixed in this variant, they introduced a flaw in the encryption method in the updated variant which we discuss below.

Master, Encrypted, Aleta, Blocking

This is very similar to the previous method, but these variants add the result of GetTickCount to the string as well, and the beginning of the user ID is a number or combination of numbers (e.g. "1", "2-4" or "201").

Example:

1-774A0371FF5858963D5C1F1A359B8A461AFB7928975491450752F4408C4AA9CB4752FCEE9-2018-01-11

The generated random string (string1) is used to derive the file encryption key.

```

000000000403688
000000000403688      loc_403688:
000000000403688      call  ?_Random_device@std@@YAIXZ ; std::_Random_device(void)
00000000040368D      8B F0      mov     esi, eax
00000000040368F      E8 CC 59 00 00      call  ?_Random_device@std@@YAIXZ ; std::_Random_device(void)
0000000004036C4      8B F8      mov     edi, eax
0000000004036C6      E8 C5 59 00 00      call  ?_Random_device@std@@YAIXZ ; std::_Random_device(void)
0000000004036CB      FF 35 FC 71 42 00      push  ds:GetClipboardOwner
0000000004036D1      0F B7 8D 8C F2 FF FF  movzx  ecx, word ptr [ebp-0D74h] ; minute
0000000004036D8      56      push  esi
0000000004036D9      FF 35 F0 71 42 00      push  ds:GetActiveWindow
0000000004036DF      89 85 C8 F1 FF FF      mov     [ebp-0E38h], eax
0000000004036E5      FF 35 4C 70 42 00      push  ds:GetCurrentThreadId
0000000004036EB      0F B7 85 8E F2 FF FF  movzx  eax, word ptr [ebp-0D72h]
0000000004036F2      51      push  ecx
0000000004036F3      50      push  eax
0000000004036F4      57      push  edi
0000000004036F5      51      push  ecx
0000000004036F6      E8 25 F7 FF FF      call  generate_rand
0000000004036FB      83 C4 04      add     esp, 4
0000000004036FE      50      push  eax
0000000004036FF      FF B5 C8 F1 FF FF      push  dword ptr [ebp-0E38h]
000000000403705      FF 15 6C 70 42 00      call  ds:GetTickCount
00000000040370B      8B 35 F8 71 42 00      mov     esi, ds:wsprintfA
000000000403711      50      push  eax
000000000403712      FF 35 5C 70 42 00      push  ds:GetCurrentProcessorNumber
000000000403718      68 88 38 43 00      push  offset aXXXXDXXXXX ; "%x%x%x%x%x%x%x%x"
00000000040371D      FF 35 78 DC 83 01      push  string1 ; LPSTR
000000000403723      FF D6      call  esi ; wsprintfA
000000000403725      0F B7 85 86 F2 FF FF  movzx  eax, word ptr [ebp-0D7Ah] ; day
00000000040372C      83 C4 34      add     esp, 34h
00000000040372F      50      push  eax
000000000403730      0F B7 85 82 F2 FF FF  movzx  eax, word ptr [ebp-0D7Eh] ; month
000000000403737      50      push  eax
000000000403738      0F B7 85 80 F2 FF FF  movzx  eax, word ptr [ebp-0D80h] ; year
00000000040373F      50      push  eax
000000000403740      FF 35 78 DC 83 01      push  string1
000000000403746      68 9C 33 43 00      push  offset unk_43339C
00000000040374B      68 D0 38 43 00      push  offset aSSD02d02d ; "%s-%s-%d-%02d-%02d"
000000000403750      68 50 B5 83 01      push  offset user_id ; LPSTR
000000000403755      FF D6      call  esi ; wsprintfA
000000000403757      B9 50 B5 83 01      mov     ecx, offset user_id
00000000040375C      C7 85 C0 F1 FF FF 00 00+mov  dword ptr [ebp-0E40h], 0
000000000403766      83 C4 1C      add     esp, 1Ch
000000000403769      C7 85 C4 F1 FF FF 0F 00+mov  dword ptr [ebp-0E3Ch], 0Fh
000000000403773      C6 85 B0 F1 FF FF 00      mov     byte ptr [ebp-0E50h], 0
00000000040377A      8D 51 01      lea    edx, [ecx+1]
00000000040377D      0F 1F 00      nop    dword ptr [eax]

```

Figure 4: Master user ID generation

In this case, the user ID generation is safe thanks to the combination of the usual random generation and the result of the GetTickCount function. However, there is another flaw in these variants, which can be used in certain cases to decrypt the encrypted files. These variants keep scanning the machine for new files to encrypt. This means that if the victim's computer hasn't been rebooted after the ransomware attack and the ransomware process hasn't been killed by the victim, then the encryption key can be retrieved from memory.

Gryphon, Crypton, Nuclear, Wyvern

These variants do not use the random generator method (with the `_srand` and `_rand` functions), which was used by all the previous variants. The user ID is generated using the result of `GetTickCount` and the address of a few functions. No other characters (like the date) are appended to it.

Example: 13828F0774A037113828F0752F440813828F013828F07549145013828F0752FCEE913828F0103EDA0

```

00000000402E9E          loc_402E9E:
00000000402E9E          add     eax, 2
00000000402E9E      83 C0 02          push   eax             ; iMaxLength
00000000402EA1      50                push   offset pszPath  ; lpString2
00000000402EA2      68 80 D2 81 01   lea   eax, [esp+4A0h+String1]
00000000402EA7      8D 84 24 90 01 00 00  push  eax             ; lpString1
00000000402EAE      50                call  ds:lstrcpynA
00000000402EAF      FF 15 60 00 41 00  call  ds:GetTickCount
00000000402EB5      FF 15 64 00 41 00  push   eax
00000000402EBB      50                push   offset sub_4026F0
00000000402EBC      68 F0 26 40 00   push   ds:GetClipboardOwner
00000000402EC1      FF 35 CC 01 41 00  push   ebx, ds:wsprintfA
00000000402EC7      8B 1D C8 01 41 00  push   offset sub_4026F0
00000000402ECD      68 F0 26 40 00   push   ds:GetCurrentThreadId
00000000402ED2      FF 35 44 00 41 00  push   offset sub_4026F0
00000000402ED8      68 F0 26 40 00   push   offset sub_4026F0
00000000402EDD      68 F0 26 40 00   push   ds:GetActiveWindow
00000000402EE2      FF 35 C0 01 41 00  push   offset sub_4026F0
00000000402EE8      68 F0 26 40 00   push   ds:GetCurrentProcessorNumber
00000000402EED      FF 35 5C 00 41 00  push   offset sub_4026F0
00000000402EF3      68 F0 26 40 00   push   offset aXXXXXXXXX ; "%XXXXXXXXXXXXXXXXXXXXXXXX"
00000000402EF8      90 96 41 00      push   string1         ; LPSTR
00000000402EFD      FF 35 B8 FC 81 01  call  ebx             ; wsprintfA
00000000402F03      FF D3            add   esp, 34h
00000000402F05      83 C4 34          push  string1
00000000402F08      FF 35 B8 FC 81 01  push  offset aS_0      ; "%s"
00000000402F0E      68 A8 96 41 00   push  offset user_id   ; LPSTR
00000000402F13      68 90 D5 81 01   call  ebx             ; wsprintfA
00000000402F18      FF D3            mov   ecx, offset user_id
00000000402F1A      B9 90 D5 81 01   mov   [esp+4A4h+var_478], 0
00000000402F1F      C7 44 24 2C 00 00 00 00  add   [esp, 0Ch]
00000000402F27      83 C4 0C          mov   [esp+498h+var_474], 0Fh
00000000402F2A      C7 44 24 24 0F 00 00 00  mov   [esp+498h+var_488], 0
00000000402F32      C6 44 24 10 00   lea  edx, [ecx+1]
00000000402F37      8D 51 01         nop
00000000402F3A      66 0F 1F 44 00 00  word ptr [eax+eax+00h]
    
```

Figure 5: Nuclear user ID generation

Payday, Shadow, Wallet

From the Payday variant, the generation of the key has changed completely. These variants use the `CryptGenRandom` function to generate 0x20 random bytes.

```

00000000404100      55                push  ebp
00000000404101      8B EC            mov   ebp, esp
00000000404103      83 EC 28        sub   esp, 28h
00000000404106      A1 08 E0 41 00  mov   eax, ___security_cookie
00000000404108      33 C5          xor   eax, ebp
0000000040410D      89 45 FC        mov   [ebp+var_4], eax
00000000404110      68 00 00 00 F0  push  0F000000h        ; dwFlags
00000000404115      6A 01          push  1                ; dwProvType
00000000404117      6A 00          push  0                ; szProvider
00000000404119      6A 00          push  0                ; szContainer
0000000040411B      8D 45 D8        lea  eax, [ebp+phProv]
0000000040411E      50                push  eax             ; phProv
0000000040411F      FF 15 20 20 41 00  call  ds:CryptAcquireContextA
00000000404125      85 C0          test  eax, eax
00000000404127      0F 84 F7 00 00 00  jz   loc_404224

0000000040412D      8D 45 DC        lea  eax, [ebp+pbBuffer]
00000000404130      50                push  eax             ; pbBuffer
00000000404131      6A 20          push  20h             ; dwLen
00000000404133      FF 75 D8        push  [ebp+phProv]    ; hProv
00000000404136      FF 15 10 20 41 00  call  ds:CryptGenRandom
0000000040413C      85 C0          test  eax, eax
0000000040413E      0F 84 D5 00 00 00  jz   loc_404219
    
```

Figure 6: Payday random generation

We can see from these examples how the developers of the BTCWare ransomware family were experimenting with the different solutions. From a completely vulnerable method they managed to reach a cryptographically secure solution. In every step, they tried to make their implementation more secure. They started with a very common mistake, using the `_srand` and `_rand` functions for random generation, which is vulnerable to brute force attack. Later, they tried to add other values to increase the randomness using `GetLocalTime`, function addresses and `GetTickCount`. By the last couple of versions, they completely removed the `_srand` and `_rand` functions and started to use the most popular and secure solution, the `CryptGenRandom` function.

User ID encryption

Once the user ID is generated, the next step is to encrypt it with the hard-coded RSA public key. This method is the same across all variants. The main functions used during encryption are the following:

`CryptStringToBinaryA`, `CryptDecodeObjectEx`, `CryptImportPublicKeyInfo`, `CryptEncrypt`.

These read, decode, and import the hard-coded RSA public key, and then encrypt the generated user ID. The encrypted user ID is saved to either a specific file, the ransom note itself, or at the end of the encrypted files. The last column of Table 2 shows this location by variants.

Symmetric key derivation

The symmetric key, used for file encryption, is derived from the user ID. The general process, which is almost exactly the same across all variants, is the following:

- The MD5 hash of the user ID (or some part of it) is calculated using the `CryptCreateHash` and `CryptHashData` functions.
- From this hash the symmetric key is derived using the `CryptDeriveKey` function. The `ALG_ID` input of this function identifies the symmetric encryption algorithm for which the key is to be generated. The fourth column of Table 2 contains the used encryption algorithm by different variants.
- Finally, if the encryption algorithm is AES, using the `CryptSetKeyParam` function the IV, padding and encryption mode are set.

The first few variants have a backup method for symmetric key generation. If for some reason the user ID generation was unsuccessful, the encryption key will be generated using the `CryptGenKey`, `CryptGetUserKey`, and `CryptExportKey` functions.

Three different encryption algorithms were used by the different variants: RC4, AES-192, and AES-256. The changes across these methods also show how the developers of the BTCWare family were experimenting with the different solutions. They started with RC4 with an implementation mistake. They use the same key for each file, which makes this method vulnerable to known plaintext attack. In the end they reached the secure AES-256 encryption in CBC-mode.

File encryption

The file encryption method is very similar for all variants. A certain number of bytes is read from the file using the `ReadFile` function. These bytes are then encrypted by `CryptEncrypt` and written to file by the `WriteFile` function. Either a new file is created for the encrypted content and the original is deleted after encryption, or the original file will be simply overwritten and renamed by `MoveFileExA`.

All the variants use specific white list of folders as mentioned below, which will be skipped during encryption.

\$recycle.bin, program files, program files (x86), programdata, windows, nvidia, intel, appdata, temp, msocache, inetpub.

The first few variants used a list of extensions to select which files to encrypt, but later variants encrypt all files regardless their extension, except in the listed folders. In most of the cases the ransom note and the already encrypted files are also skipped, but we have seen a few exceptions.

For example, one of the Aleta samples encrypted its own ransom note, which makes it impossible for the victims to read the attackers' message.

Some variants also have a list of folders where they do encrypt despite of the previous exclusion:

- C:\Program Files\MySQL\
- C:\Program Files (x86)\Firebird\
- C:\Program Files (x86)\MSSQL.1\

It's not just the flaws in the implementation of the cryptographic solutions that made it possible to decrypt certain variants of this ransomware family. The developers of BTCWare have made two of their master secret keys public.

The first release made possible to decrypt the BTCWare, Cryptobyte, Cryptowin variants, while the second key belongs to the Theva, Onyon, Xfile, and Master variants.⁴⁵ (There are some exceptions among these variants which use these extensions, but a different private key.) Every time the key was released a new variant appeared with a different private key to keep the business running.

Ransom demand

The ransom note is almost exactly the same for all variants. The file format and design changed a little bit over time, but the main content remained the same. The victims have to contact the attackers using a given email address. They can send a couple of files to the attackers who will decrypt those for free.

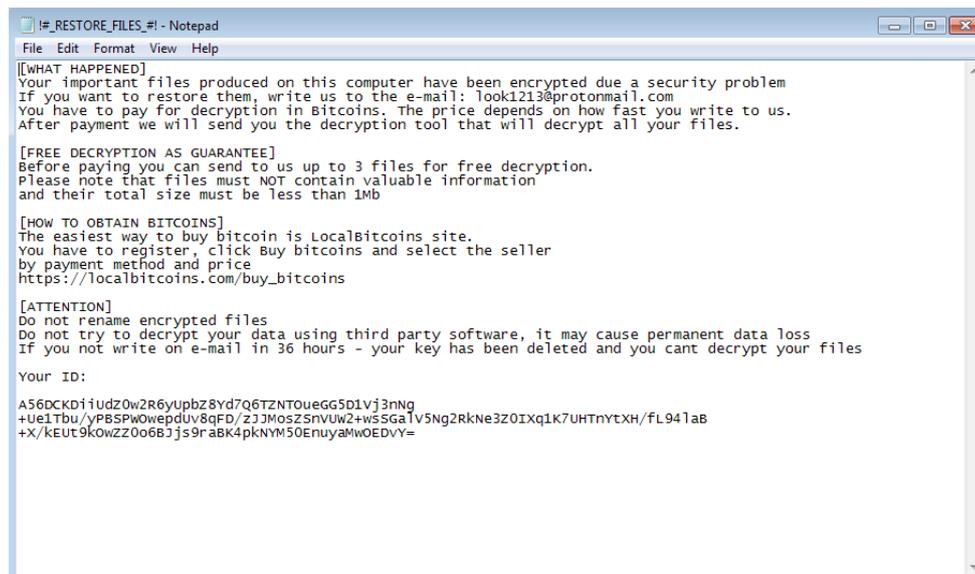


Figure 7: Master ransom note

In a few cases we have seen different ransom notes but with much less content.

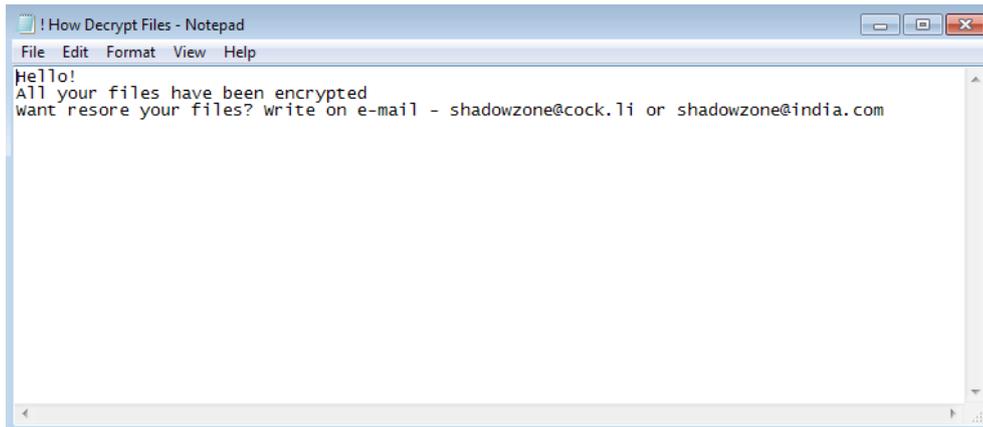


Figure 8: Wallet ransom note



Figure 9: PayDay ransom note

We have found a few examples in which the attackers provided a TOR address to the victims instead of an email address. For example, the address `hxxp://cr7icbfqm64hixta[.]onion` is used for payment instead of contacting the author via email ID.

Ransom payment

If a victim contacts the attackers using the given email address, the victim will receive a Bitcoin address and the amount of Bitcoin demanded. We tried to collect Bitcoin addresses, mainly from forums, where the victims shared the information they received from the attackers. Following one Bitcoin address we found on the Bleeping Computer support forum, we reached a Bitcoin address in two steps with more than 925 Bitcoins.

Bitcoin Address

Addresses are identifiers which you use to send bitcoins to another person.

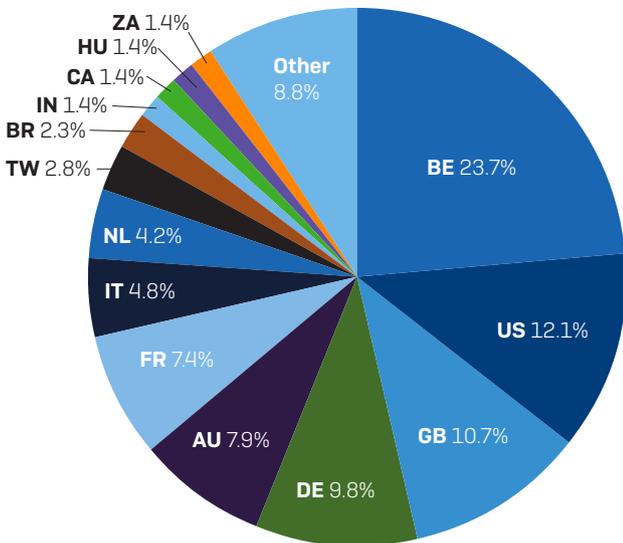
Summary	Transactions	
Address: 1M1Pc1SVHUJGiiScvcX7VWjx82nUDnniwM	No. Transactions: 16	
Hash 160: db75625a79254f74e3cea358dd12e19e695cd933	Total Received: 925.51341941 BTC	
Tools: Related Tags - Unspent Outputs	Final Balance: 925.51341941 BTC	

Request Payment
Donation Button

Statistics

Although most of the BTCWare attacks targeted European countries with Belgium in the lead, we have seen attacks all around the world. The United States was the second most affected country and we can find Taiwan, India, Canada, and the South African Republic represented among the targets as well.

In the second half of 2017 BTCWare has shown a steady presence, with relatively low number of customer lookups, thanks to the fact that it was distributed using RDP.



Protection

Sophos detects BTCWare ransomware using the following detections:

Mal/Btcware-A, Troj/Btcware-*, HPmal/BTCWare-A.

Additionally, Sophos [Intercept X](#) proactively prevents the malware from attacking your data, as the CryptoGuard component stops the ransomware from scrambling your files.

Summary

Ransomware in general is a big problem and users need to apply good security practices to avoid falling victim to it.

For more information about ransomware and best security practices, read [“How to stay protected against ransomware.”](#)

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Appendix A

Here are some of the email addresses to contact the attackers mentioned in the ransom notes:

3bitcoins[at]protonmail.com	kekin[at]cock.li
alekstraza[at]bigmir.net	keyforyou[at]tuta.io
assistance[at]firemail.cc	lavandos[at]dr.com
atf13[at]tuta.io	lockers[at]tutamail.com
avalona.toga[at]aol.com	lockers[at]protonmail.com
averia[at]tuta.io	mail[at]aleta.cc
averiasw[at]qq.com	Merlin[at]aolonline.top
badking[at]india.com	mia.kokers[at]aol.com
barbarosso1051223[at]tutanota.com	microcost[at]bigmir.net
black.block[at]qq.com	moneymaker2[at]india.com
black.mirror[at]qq.com	mortalis_certamen[at]aol.com
black.world[at]tuta.io	nicecrypt[at]india.com
blacklandfat[at]qq.com	nkr.siger[at]protonmail.ch
bravobravo[at]cock.li	nuclear[at]cryptmaster.info
Checkzip[at]india.com	pakhomovsemen60[at]gmail.com
chines34[at]protonmail.ch	payfordecrypt[at]qq.com
crypt24[at]protonmail.com	predatorthre[at]bigmir.net
darkwaiderr[at]cock.li	prt.nyke[at]protonmail.ch
darkwaiderr[at]tutanota.com	realsapport[at]bigmir.net
decr[at]cock.li	roggers[at]bigmir.net
decryptorx[at]cock.li	stopstorage[at]qq.com
decryptyourfilesheeee1[at]cock.li	support[at]fbamasters.com
garryhelpyou[at]qq.com	torettoxyx[at]gmail.com
gladius_rectus[at]aol.com	tsce308[at]tutanota.com
goliath[at]tuta.io	tullakump[at]tutanota.com
help[at]jonyon.info	universe1[at]protonmail.ch
hostname[at]bigmir.net	usermsd[at]cock.li
info[at]kraken.cc	vargbt[at]tutanota.com
info[at]zayka.pro	webmafia[at]asia.com
irmagetstein[at]india.com	wyvern[at]cryptmaster.info

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